COUNTING ARGUMENTS*

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One may have experienced in everyday life that a problem which seems to be very difficult can be solved easily by changing one's point of view.

This is also true in mathematics and flexibility is one of the most important factors in getting results in mathematics.

Notation. |X| or #X denotes the number of elements of a set X, and $X^{(k)}$ denotes the family of k-subsets of X, where a k-subset means a subset of k elements.

Example 1. Let N be a set of elements. Then

 $|\mathbb{N}^{(k)}| = \left(\begin{array}{c} n\\ k \end{array}\right).$

Count the number of elements of the set

 $X = \{ (i, S) \in N \times N^{(k)} \mid i \in S \}$

in two ways. There are n choices of i, and for each i the

number of k-subsets containing i is $\begin{pmatrix} n-1\\ k-1 \end{pmatrix}$. Hence

 $|X| = n \left(\begin{array}{c} n - 1 \\ k - 1 \end{array} \right),$

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On the other hand, there are $\begin{bmatrix} n \\ k \end{bmatrix}$ choices of S, and for each S the number of $i \in S$ is k. Hence $|X| = k \begin{bmatrix} n \\ k \end{bmatrix}$, and we have

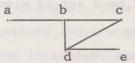
$$\left(\begin{array}{c}n\\k\end{array}\right) = \frac{n}{k} \left(\begin{array}{c}n-1\\k-1\end{array}\right).$$

Then by induction, we have

$$\left(\begin{array}{c}n\\k\end{array}\right)=\frac{n(n-1)\ldots(n-k+1)}{k(k-1)\ldots 1}.$$

Example 2. Let V be a set of v elements and E a subset of $V^{(2)}$. Then we call $\Gamma = (V, E)$ a graph and the elements of V and E are called vertices and edges respectively. For $a \in V$, the number of edges containing a is called the *degree* of a and is denoted by d(a).

For instance, the following figure



gives a graph $\Gamma = (V, E)$ with $V = \{a, b, c, d, e\}$ and

 $E = \{\{a, b\}, \{b, c\}, \{b, d\}, \{c, d\}, \{d, e\}\},\$

and d(a) = 1, d(b) = 3, d(c) = 2, d(d) = 3, d(e) = 1.

xEV

Now for a graph $\Gamma = (V, E)$ count the number of elements of the set

 $X = \{(x, e) \in V \times E \mid x \in e \}$

in two ways. Then we have

$$\Sigma d(\mathbf{x}) = 2|\mathbf{E}|,$$

and the following holds.

Theorem. $\Sigma \ d(x)$ is even. $x \in V$ As a corollary, we have the following Corollary. (1) #{ $x \in V \mid d(x)$ is odd } is even. (2) If |V| is odd then there is an $x \in V$ with d(x) even. In mathematics it happens very often that one thing assumes two or more different aspects. As an example, I will talk about the Marriage Theorem and Dilworth's Theorem.

Example 3. Let G be a set of m girls, B a set of n boys and let

 $F = \{ (g, b) \in G \times B \mid b \text{ is a boy friend of } g \}.$ For $g \in G$ denote $\{ b \in B \mid (g, b) \in F \}$ by B(g). Then the question is: under what conditions can every girl choose her husband from her boy friends? The answer is given in the following theorem.

Theorem 1. Every girl can choose her husband from her boy friends if and only if for every $k \le m$ and for every k girls g_1, \ldots, g_k ,

 $\left|\bigcup_{i=1}^{k} B(g_{i})\right| \leq k.$

Proof. Necessity is obvious.

Sufficiency can be proved by using Dilworth's Theorem.

Let P be a set with order \leq , and assume that the order satisfies the following axioms:

(1) $a \leq a$ for every $a \in P$,

(2) $a \le b$, $b \le a \Rightarrow a = b$ for every $a, b \in P$,

(3) $a \le b$, $b \le c \Rightarrow a \le c$ for every $a, b, c \in P$.

Then we call P a partially ordered set, where "partially" means that some two elements a and b may not be related with respect to \leq and in this case we say that a and b are *independent*. More generally, a_1, \ldots, a_r are *independent* if every two elements a_i , a_j ($i \neq j$) are independent. We denote the maximum of the number of independent elements of P by M(P). On the other hand, a subset C of P is called a *chain* if for any two elements a, b of C either $a \le b$ or else $b \le a$. If $P = C_1 \cup \ldots \cup C_s$ (disjoint union) where the C_i are chains, then this union is called a *chain decomposition* of P. Let m(P) be the maximum of the number of chains appearing in a chain decomposition of P.

Theorem 2 (Dilworth). M(P) = m(P).

Theorem 2 implies Theorem 1.

Proof. Let $P = G \cup B$ and define an order \leq on P as follows:

 $x \le y \Leftrightarrow y = x \text{ or } (y, x) \in F.$

Then a chain is .x or $\begin{bmatrix} g \\ b \end{bmatrix}$. Let M(P) = m(P) = M. Then there

is a chain decomposition

and M = n + (m - r). We will show that m - r = 0. By Theorem 2, there is an independent set

$$g_1 \cdot \cdot \cdot g_s$$

 $b_1 \cdot \cdot \cdot b_t$

with s + t = M = n + (m - r). Then

$$\{b_1, \ldots, b_t\} \subseteq B - \bigcup_{i=1}^{U} B(g_i),$$

and we have $t \le n - s$. Thus $s + t = n + (m - r) \le n$ and we have m - r = 0.